

# Designing National IP/MPLS Networks with Flexgrid Optical Technology

L. Velasco<sup>1\*</sup>, P. Wright<sup>2</sup>, A. Lord<sup>2</sup>, G. Junyent<sup>1</sup>

<sup>(1)</sup> Universitat Politècnica de Catalunya (UPC), Barcelona (Spain), Email: lvelasco@ac.upc.edu

<sup>(2)</sup> British Telecom Innovate & Design, Ipswich (United Kingdom)

**Abstract** We propose a two-step procedure to design flexgrid-based national networks. Locations are first partitioned into a set of metro areas interconnected through a flexgrid optical network. Next, each network is designed separately. Optimal results show a future large (>200 nodes) flexgrid core network inter-connecting small (~10 nodes) metro regions.

## Introduction

The problem of designing IP/MPLS networks consists of finding the configurations of the whole set of routers and links so as to transport a given traffic matrix whilst minimizing capital expenditure (CAPEX). As a consequence of link lengths, national IP/MPLS networks have been built on the top of optical networks, and thus the design problem has been typically addressed through a multilayer IP/MPLS-over-optical approach where IP/MPLS routers are placed side-by-side of optical cross-connect (OXC) <sup>1</sup>. Besides, due to the coarse granularity of the current fixed-grid wavelength division multiplexing technology used at the optical layer, aggregation networks are deployed between clients and core networks to groom client flows thus increasing the spectral efficiency of the optical layer.

However, flexgrid technology <sup>2</sup>, providing a finer optical spectrum granularity, allows a new approach <sup>3,4</sup>: IP/MPLS routers equipped with multi-flow transponders (MF-TPs) <sup>5</sup> are connected to bandwidth-variable (BV) OXC <sup>6</sup>, transforming the multilayer approach into a single-layer approach where a number of IP/MPLS metro area networks performing aggregation are connected through a flexgrid-based core network (Fig. 1).

In this paper we propose a two-step procedure to design such networks and investigate the optimal size of those aggregation metro areas so as to deploy spectrum-efficient flexgrid-based core optical networks.

## Network design proposed procedure

Given a national network, such as BT's, with over 1000 node locations, a full, single step network optimization is clearly impossible. In this paper we present a two-step network design procedure which takes advantage of the separation model described above:

1. The network is partitioned into metro areas, i.e. locations are grouped into a number of metro areas. Internal metro area traffic

between locations in the same area and aggregated traffic between metro areas is computed. Aggregation traffic will be conveyed by the optical core network.

2. IP/MPLS metro and optical core networks are designed, taking as input each traffic matrix computed in 1) and a network topology.

Fig. 2 illustrates the proposed design procedure where some data needs to be pre-computed: candidates for metro nodes can be limited to those nodes with sufficient connectivity degree and each location can have a limited subset of potential metro nodes (e.g. based on distance). In this paper we focus on the first step.

## Metro Area Partitioning (MAP) Problem

The MAP problem targets at maximizing the total aggregated traffic to be conveyed by the core network. Note that by maximizing the aggregated traffic we are indirectly minimizing the internal traffic of the resulting metro areas. That is important since the cost per bit of the optical switching technology is generally cheaper than that of the IP/MPLS and the aggregated flows in the core network occupy many spectral resources, so are most cost effectively groomed and routed optically.

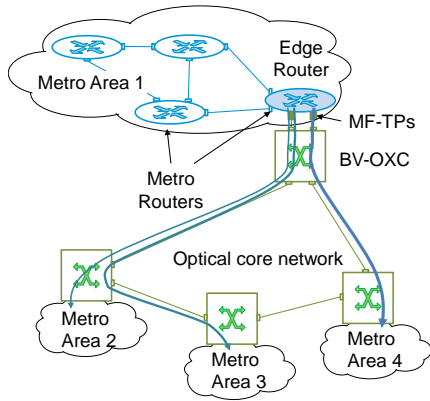
The MAP problem can be formally stated as:

**Given:** *i*) a set  $L$  of Locations, *ii*) a subset  $M \subseteq L$  of locations that can be selected as metro locations, *iii*) the subset  $M(l) \subseteq M$  of metro locations where each location  $l$  can belong. *iv*) the location-to-location traffic matrix, and *v*) the considered slot width ( $\Delta f$ ) for the flexgrid-based core network.

**Output:** The set of metro areas. For each metro area  $m$ , the locations belonging to that area and its internal traffic. The set of aggregated traffic between metro areas ( $b_{mm'}$ ).

**Objective:** Maximize the aggregated traffic exchanged between metro areas ( $\sum_{m \in M} \sum_{m' \in M} b_{mm'}$ ) subject to ensuring a minimum spectral efficiency for the core network.

Note that spectral efficiency links the two proposed steps so that the final network design



**Fig. 1** Network design: metro areas are equipped with IP/MPLS routers. The core network is based on the flexgrid technology.

eventually minimizes CAPEX. It is defined as follows, where  $B_{mod}$  represents the spectral efficiency of the chosen modulation format.

$$\sum_{m \in M} \sum_{\substack{m' \in M \\ m' \neq m}} \frac{b_{mm'}}{\Delta f \cdot B_{mod}} \bigg/ \sum_{m \in M} \sum_{\substack{m' \in M \\ m' \neq m}} \left\lceil \frac{b_{mm'}}{\Delta f \cdot B_{mod}} \right\rceil \quad (1)$$

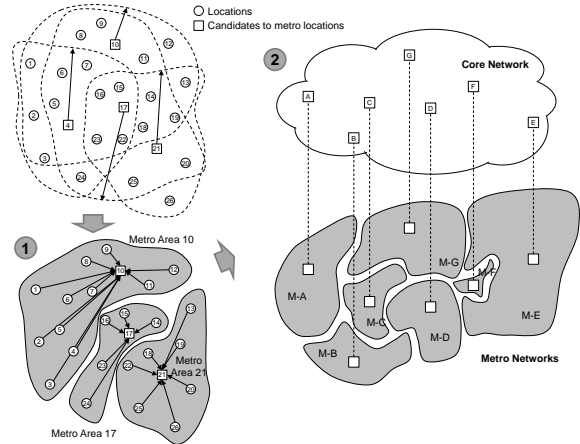
Note that the ceiling operation computes the amount of slots to convey the requested data flow under the chosen slot width.

The MAP problem was modeled using mixed integer linear programming (MILP). However, its exact solution becomes impractical when realized network and traffic instances are considered. Hence, we developed a heuristic algorithm. The performance of the heuristic was compared against the optimal solution obtained solving the MILP model over small instances where the optimal solution was found in all the tests performed.

### Illustrative numerical results

In this section, we present the obtained results from solving a close-to-real problem instance consisting of 1113 locations, based on the BT network. Those locations (323) with a connectivity degree of 4 or above were selected as potential metro locations. A 3.22 Pb/s traffic matrix was obtained by considering the number of residential and business premises in the proximity of each location. Locations could only be parented to a potential metro node if they were within a 100 km radius. Each result reported in the following is the average of four different executions of the heuristic algorithm, each running 48 hours.

Fig. 3 plots the amount of aggregated traffic for each solution from solving the heuristic algorithm when the number of metro areas is fixed. The relationship between the amount of aggregated traffic and the number of metro areas is clearly shown: more metro areas entail higher aggregated traffic to be exchanged. On the contrary, it is worth noting that the amount of



**Fig. 2** Two-step procedure for network design.

aggregated traffic does not depend on the slot width. As a result of the problem's objective, i.e. maximizing the amount of aggregated traffic, its increment when all the 323 metro areas are opened is only 6% with respect to just 20 areas. In order to understand the behavior of the core network spectral efficiency, this has been included as part of the objective function to be maximized. Therefore, plots in Fig. 3 show the maximum spectral efficiency of the solutions as a function of the slot width and the number of metro areas selected. As illustrated, network spectral efficiency decreases sharply when the number of metro areas is increased since more flows with lower amount of traffic are needed to be transported over the core network. Note that the traffic matrix to be transported by the core network has  $|M| \cdot (|M| - 1)$  unidirectional flows. Let us consider a threshold for spectral efficiency of 80% (horizontal red line in Fig. 3). Then, the largest number of metro areas is 116, 165, 216, or 295 when the 50, 25, 12.5 and 6.25 GHz slot width, respectively, is selected. Obviously, the coarser the grid granularity chosen for the optical network, the larger the metro areas have to be for the spectral efficiency threshold selected and thus the lower the number of metro areas opened.

Fig. 4 presents the details of the solutions as functions of the number of metro areas, while Table 1 focuses on the characteristics of those solutions in the defined spectral efficiency threshold for each slot width. Note that plots in Fig. 4, as for Fig. 3, are valid irrespective of the slot width used, since no spectral efficiency threshold was required.

In Fig. 4a, the size of the metro areas is represented. Note that in the reference interval [116, 295] the average number of locations in each metro area is lower than 10, being lower than 20 for the largest metro area. This limited metro area size simplifies the design of those

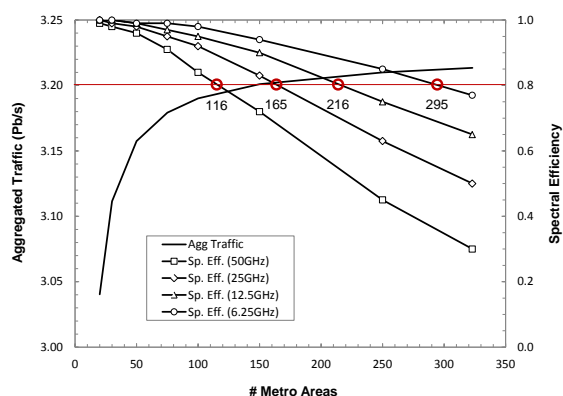


Fig. 3 Aggregated traffic and spectral efficiency.

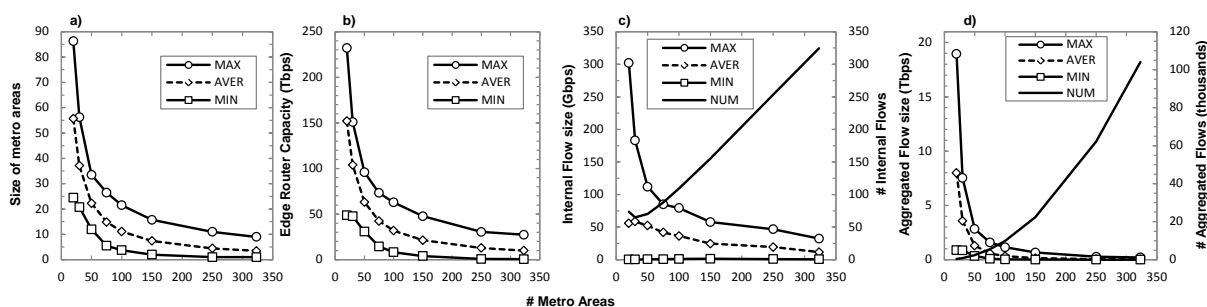


Fig. 4 Details of the solutions against the number of metro areas. Max, min, and average plots are shown.

IP/MPLS networks. Besides, switching capacity of metro edge IP/MPLS routers is proportional to the size of the metro areas as shown in Fig. 4b. They require capacities of up to 58 Tb/s (28.5 Tb/s on average) when the 50GHz slot width is used, decreasing to 28.5 Tb/s (11 Tb/s on average) for 6.25GHz. Hence, finer slot widths might keep routers to single chassis size - far more efficient and cost effective.

Fig. 4c surveys the size of the metro area internal data flows (flows where at least one end router is in a given metro area). In the reference interval, the size of the internal flows is up to 72 Gb/s, 32 Gb/s on average, when the 50GHz slot width is used, decreasing to 38 Gb/s, 15 Gb/s on average, when the 6.25 GHz slot width is used. The average size of the metros traffic matrix is small as a result of the size of the metro areas: lower than 300 unidirectional flows. Finally, Fig. 4d examines the size of the aggregated data flows in the core network. When the number of metro areas is low, the size of the aggregated flows is quite large but the overall size of the traffic matrix is quite small: the largest flow conveys 19 Tb/s although there are only 380 aggregated data flows. In comparison, when the number of metro areas is large, the size of the largest flow is smaller but the size of the core network traffic matrix is very large: the largest flow conveys 203 Gb/s but there are as many as 104,000 aggregated data flows.

In the reference interval (see Table 1), the size of the aggregated flows is up to 989 Gb/s, 264

Table 1. Solutions details (spectral efficiency = 0.8)

Slot width	50 GHz	25 GHz	12.5 GHz	6.25 GHz
# Metro Areas	116	165	216	295
Metro areas size (Max / AVer)	19.66 / 9.94	15.04 / 6.97	15.04 / 5.46	9.77 / 3.83
Edge Router Capacity (Max / AVer (Tb/s))	58.0 / 28.5	45.1 / 20.1	35.3 / 15.7	28.5 / 11.1
Metro flows (Max / AVer (Gb/s) / #)	72.33 / 32.68 / 124	56.06 / 23.84 / 170	50.52 / 21.12 / 220	37.92 / 14.60 / 297
Agg. flows (Max / AVer (Gb/s) / #)	989.06 / 264.95 / 13,884	602.16 / 129.47 / 28,335	404.51 / 82.73 / 48,684	229.91 / 38.83 / 87,990

Gb/s on average, when the 50GHz slot width is used, decreasing to 223 Gb/s, 39 Gb/s on average, when the 6.25 GHz slot width is used. The size of the optical network traffic matrix increases from 14,000 to 48,000 unidirectional flows when the 50 and 12.5 GHz slots widths are selected.

## Conclusions

A two-step procedure to design flexgrid-based national networks consisting of a set of metro areas interconnected through a flexgrid optical core has been proposed. The MAP problem to partition the network into metro areas has been stated and solved for the BT network. Results showed that simpler and smaller metro areas containing 10-15 locations are enough to obtain good spectral efficiency in the flexgrid-base core network when 12.5 or even 6.25 GHz slot widths are used.

## Acknowledgements

The research leading to these results has received funding from the European Community's Seventh Framework Programme FP7/2007-2013 under grant agreement no. 247674 STRONGEST project and by the MINECO project ELASTIC (TEC2011-27310).

## References

- [1] M. Ruiz, et al., J. Opt. Comm. Netw., **3** (2011)
- [2] M. Jinno et al., IEEE Commun. Mag., **47**, (2009)
- [3] O. Gerstel, Proc. Photonics in Switching (2010)
- [4] M. Jinno et al., Proc. ECOC (2011)
- [5] H. Takara et al., Proc. ECOC (2011)
- [6] G. Wellbrock, Proc. OFC/NFOEC (2011)